





### Sparse Days 2025

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# Mixed precision preconditioned GMRES

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Alfredo Buttari, Nicholas J. Higham, T. M., and Bastien Vieublé. A modular framework for the backward error analysis of GMRES. HAL EPrint hal-04525918, 2024, to appear in IMAJNA. https://hal.science/hal-04525918

Alfredo Buttari, Xin Liu, T. M., and Bastien Vieublé. Mixed precision strategies for preconditioned GMRES: a comprehensive analysis. HAL EPrint hal-05071696, 2025. https://hal.science/hal-05071696

# GMRES

- Krylov-based iterative solver for the solution of general square linear systems Ax = b
- Compute an orthonormal Krylov basis V<sub>k</sub> through the Arnoldi iteration
- Find the vector xk in span{Vk} that minimizes ||Axk − b||
- Reiterate until x<sub>k</sub> is a satisfactory approximation of x

### GMRES

**Input:**  $A \in \mathbb{R}^{n \times n}$ ,  $b \in \mathbb{R}^n$ ,  $x_0 \in \mathbb{R}^n$ .  $\tau > 0$ 1:  $r_0 = b - Ax_0$ 2:  $\beta = ||r_0||, v_1 = r_0/\beta, k = 1$ 3: repeat 4:  $W_{k} = A_{V_{k}}$ 5: **for** i = 1, ..., k **do** 6:  $h_{i,k} = v_i^T w_k$ 7:  $w_k = w_k - h_{i,k} v_i$ 8: end for 9:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 10:  $V_k = [v_1, \ldots, v_k]$ 11:  $H_k = \{h_{i,i}\}_{1 \le i \le i+1: 1 \le i \le k}$ 12:  $y_k = \operatorname{argmin}_{v} \|\beta e_1 - H_k y\|$ 13: k = k + 114: until  $\|\beta e_1 - H_k v_k\| < \tau$ 15:  $x_{k} = x_{0} + V_{k} v_{k}$ **Output:**  $x_k \in \mathbb{R}^n$  such that  $||Ax_k - b|| \leq \tau$ 

### Preconditioned GMRES

$$M^{-1}Ax = M^{-1}b$$

#### Left GMRES

1:  $r_0 = b - Ax_0$ 2:  $s_0 = M^{-1}r_0$ 3:  $\beta = ||s_0||, v_1 = s_0/\beta, k = 1$ 4: repeat 5:  $z_k = Av_k$ 6:  $w_{k} = M^{-1}z_{k}$ 7: for i = 1, ..., k do 8:  $h_{i,k} = v_i^T w_k$ 9:  $w_k = w_k - h_{i,k} v_i$ 10: end for 11:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 12:  $V_k = [v_1, \ldots, v_k]$ 13:  $H_k = \{h_{i,i}\}_{1 \le i \le i+1:1 \le i \le k}$ 14:  $y_k = \operatorname{argmin}_{v} \|\beta e_1 - H_k y\|$ 15: k = k + 116: until  $\|\beta e_1 - H_k y_k\| < \tau$ 17: 18:  $x_k = x_0 + V_k y_k$ 

$$AM^{-1}y = b$$
,  $x = M^{-1}y$ 

#### **Right GMRES**

1:  $r_0 = b - Ax_0$ 2: 3:  $\beta = ||r_0||, v_1 = r_0/\beta, k = 1$ 4: repeat 5:  $z_{\mu} = M^{-1} v_{\mu}$ 6:  $w_k = A z_k$ 7: for i = 1, ..., k do 8:  $h_{i,k} = v_i^T w_k$ 9:  $w_k = w_k - h_i k v_i$ 10: end for 11:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 12:  $V_k = [v_1, \ldots, v_k]$ 13:  $H_k = \{h_{i,i}\}_{1 \le i \le i+1; 1 \le j \le k}$ 14:  $y_k = \operatorname{argmin}_v \|\beta e_1 - H_k y\|$ 15: k = k + 116: until  $\|\beta e_1 - H_k y_k\| < \tau$ 17:  $d_k = V_k y_k$ 18:  $x_k = x_0 + M^{-1} d_k$ 

### Preconditioned GMRES

$$M^{-1}Ax = M^{-1}b$$

#### Left GMRES

1:  $r_0 = b - Ax_0$ 2:  $s_0 = M^{-1}r_0$ 3:  $\beta = ||s_0||, v_1 = s_0/\beta, k = 1$ 4: repeat 5:  $z_k = Av_k$ 6:  $w_{k} = M^{-1}z_{k}$ 7: for i = 1, ..., k do 8:  $h_{i,\nu} = v_{i}^{T} w_{\nu}$ 9:  $w_k = w_k - h_{i,k} v_i$ 10: end for 11:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 12:  $V_k = [v_1, \ldots, v_k]$ 13:  $H_k = \{h_{i,i}\}_{1 \le i \le i+1:1 \le i \le k}$ 14:  $y_k = \operatorname{argmin}_{v} \|\beta e_1 - H_k y\|$ 15: k = k + 116: until  $\|\beta e_1 - H_k y_k\| < \tau$ 17: 18:  $x_k = x_0 + V_k y_k$ 

$$AM^{-1}y = b$$
,  $x = M^{-1}y$ 

#### Flexible GMRES

1:  $r_0 = b - Ax_0$ 2: 3:  $\beta = ||r_0||, v_1 = r_0/\beta, k = 1$ 4: repeat 5:  $z_{\mu} = M^{-1} v_{\mu}$ 6:  $w_k = A z_k$ 7: for i = 1, ..., k do 8:  $h_{i,k} = v_i^T w_k$ 9:  $w_k = w_k - h_{i,k} v_i$ 10: end for 11:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 12:  $V_k = [v_1, \ldots, v_k], Z_k = [z_1, \ldots, z_k]$ 13:  $H_k = \{h_{i,j}\}_{1 \le i \le i+1:1 \le i \le k}$ 14:  $y_k = \operatorname{argmin}_v \|\beta e_1 - H_k y\|$ 15: k = k + 116: until  $\|\beta e_1 - H_k v_k\| < \tau$ 17: 18:  $x_k = x_0 + Z_k d_k$ 

# Mixed precision

	Signif. bits	Exp. bits	Range ( $f_{\rm max}/f_{\rm min}$ )	Unit roundoff <i>u</i>
fp128	113	15	$2^{32766}pprox 10^{9863}$	$2^{-114}\approx1\times10^{-34}$
fp64	52	11	$2^{2046}pprox 10^{308}$	$2^{-53}pprox 1 imes 10^{-16}$
fp32	23	8	$2^{254}pprox 10^{76}$	$2^{-24}pprox 6 imes 10^{-8}$
tfloat32	10	8	$2^{254}pprox 10^{76}$	$2^{-11}pprox 5 imes 10^{-4}$
fp16	10	5	$2^{30}pprox 10^9$	$2^{-11}\approx5\times10^{-4}$
bfloat16	7	8	$2^{254}pprox 10^{76}$	$2^{-8}pprox 4 imes 10^{-3}$
fp8 (E4M3)	3	4	$2^{15}pprox 3 imes 10^4$	$2^{-4}pprox 6 imes 10^{-2}$
fp8 (E5M2)	2	5	$2^{30}pprox 10^9$	$2^{-3}pprox 1 imes 10^{-1}$
fp6 (E2M3)	3	2	$2^3 \approx 8$	$2^{-4}pprox 6 imes 10^{-2}$
fp6 (E3M2)	2	3	$2^7 pprox 128$	$2^{-3}pprox 0.125$
fp4 (E2M1)	1	2	$2^3 \approx 8$	$2^{-2} pprox 0.25$

Lower precisions:

- © Faster, consume less memory and energy
- Cover accuracy and narrower range
- $\Rightarrow$  Mixed precision algorithms

■ N. J. Higham and T. M. Mixed precision algorithms in numerical linear algebra. Acta Numerica 2022. https://www.doi.org/10.1017/S0962492922000022

### Left GMRES

1:	$r_0 = b - Ax_0$	ua
2:	$s_0 = M^{-1} r_0$	U <sub>m</sub>
3:	$eta = \  s_0 \ , \; v_1 = s_0 / eta, \; k = 1$	ug
4:	repeat	
5:	$z_k = A v_k$	u <sub>a</sub>
6:	$w_k = M^{-1} z_k$	U <sub>m</sub>
7:	for $i = 1,, k$ do	
8:	$h_{i,k} = v_i^T w_k$	ug
9:	$w_k = w_k - h_{i,k} v_i$	ug
10:	end for	
11:	$h_{k+1,k} = \ w_k\ , v_{k+1} = w_k/h_{k+1,k}$	ug
12:	$V_k = [v_1, \ldots, v_k]$	
13:	$H_k = \{h_{i,j}\}_{1 \le i \le j+1; 1 \le j \le k}$	
14:	$y_k = \operatorname{argmin}_{v} \ \beta e_1 - H_k y\ $	ug
15:	k = k + 1	
16:	until $\ \beta e_1 - H_k y_k\  \leq \tau$	
17:		
18:	$x_k = x_0 + V_k y_k$	ug

### **Right GMRES**

1:	$r_0 = b - A x_0$	ua
2:		
3:	$\beta =   r_0  , \ v_1 = r_0/\beta, \ k = 1$	ug
4:	repeat	
5:	$z_k = M^{-1} v_k$	u <sub>m</sub>
6:	$w_k = A z_k$	Ua a
7:	for $i = 1,, k$ do	
8:	$h_{i,k} = v_i^T w_k$	u <sub>g</sub>
9:	$w_k = w_k - h_{i,k} v_i$	u <sub>g</sub>
10:	end for	
11:	$h_{k+1,k} = \ w_k\ , v_{k+1} = w_k/h_{k+1,k}$	u <sub>g</sub>
12:	$V_k = [v_1, \ldots, v_k]$	
13:	$H_k = \{h_{i,j}\}_{1 \le i \le j+1; 1 \le j \le k}$	
14:	$y_k = \operatorname{argmin}_{y} \ eta e_1 - H_k y\ $	ug
15:	k = k + 1	
16:	until $\ eta e_1 - H_k y_k\  \leq  au$	
17:	$d_k = V_k y_k$	ug
18:	$x_k = x_0 + M^{-1}d_k$	u <sub>m</sub>

### Left GMRES

1:	$r_0 = b - Ax_0$	ua
2:	$s_0 = M^{-1} r_0$	U <sub>m</sub>
3:	$eta = \  s_0 \ , \; v_1 = s_0 / eta, \; k = 1$	ug
4:	repeat	
5:	$z_k = A v_k$	u <sub>a</sub>
6:	$w_k = M^{-1} z_k$	U <sub>m</sub>
7:	for $i = 1,, k$ do	
8:	$h_{i,k} = v_i^T w_k$	ug
9:	$w_k = w_k - h_{i,k}v_i$	ug
10:	end for	
11:	$h_{k+1,k} = \ w_k\ , v_{k+1} = w_k/h_{k+1,k}$	ug
12:	$V_k = [v_1, \ldots, v_k]$	
13:	$H_k = \{h_{i,j}\}_{1 \le i \le j+1; 1 \le j \le k}$	
14:	$y_k = \operatorname{argmin}_{v} \ \beta e_1 - H_k y\ $	ug
15:	k = k + 1	
16:	until $\ eta e_1 - H_k y_k\  \leq  au$	
17:		
18:	$x_k = x_0 + V_k y_k$	ug

#### Flexible GMRES

1:	$r_0 = b - A x_0$	ua	
2:			
3:	$\beta =   r_0  , \ v_1 = r_0/\beta, \ k = 1$	ug	
4:	repeat	-	
5:	$z_k = M^{-1} v_k$	u <sub>m</sub>	
6:	$w_k = A z_k$	<b>U</b> a	
7:	for $i = 1,, k$ do		
8:	$h_{i,k} = v_i^T w_k$	ug	
9:	$w_k = w_k - h_{i,k} v_i$	ug	
10:	end for		
11:	$h_{k+1,k} = \ w_k\ , v_{k+1} = w_k/h_{k+1,k}$	ug	
12:	$V_k = [v_1, \ldots, v_k], \ Z_k = [z_1, \ldots, z_k]$		
13:	$H_k = \{h_{i,j}\}_{1 \le i \le j+1; 1 \le j \le k}$		
14:	$y_k = \operatorname{argmin}_{y} \ eta e_1 - H_k y\ $	ug	
15:	k = k + 1		
16:	until $\ eta e_1 - H_k y_k\  \leq  au$		
17:			
18:	$x_k = x_0 + Z_k d_k$	um	

### Mixed precision GMRES: state-of-the-art

- $u_a = u_m \ll u_g$  Applying A and  $M^{-1}$  in high precision to improve accuracy. Existing studies are dedicated to left-preconditioned GMRES
  - E. Carson and N. J. Higham. A New Analysis of Iterative Refinement and Its Application to Accurate Solution of III-Conditioned Sparse Linear Systems. SISC 2017.
  - E. Carson and N. J. Higham. Accelerating the solution of linear systems by iterative refinement in three precisions. SISC 2018.
  - E. Carson, N. J. Higham, and S. Pranesh. Three-Precision GMRES-Based Iterative Refinement for Least Squares Problems. SISC 2020.
  - P. R. Amestoy, A. Buttari, N. J. Higham, J.-Y. L'Excellent, T. M., and B. Vieublé. Five-Precision GMRES-Based Iterative Refinement. SIMAX 2024.
  - P. R. Amestoy, A. Buttari, N. J. Higham, J.-Y. L'Excellent, T. M., and B. Vieublé. Combining sparse approximate factorizations with mixed-precision iterative refinement. TOMS 2023.
  - E. Carson and N. Khan. Mixed Precision Iterative Refinement with Sparse Approximate Inverse Preconditioning. SISC 2023.

### Mixed precision GMRES: state-of-the-art

- $u_a = u_m \ll u_g$  Applying A and  $M^{-1}$  in high precision to improve accuracy. Existing studies are dedicated to left-preconditioned GMRES
- $u_a = u_g \ll u_m$  Applying  $M^{-1}$  in a lower precision to improve performance. Existing studies are dedicated to flexible GMRES

🖹 M. Arioli and I. S. Duff. Using FGMRES to obtain backward stability in mixed precision. ETNA 2008.

J. D. Hogg and J. A. Scott. A fast and robust mixed-precision solver for the solution of sparse symmetric linear systems. TOMS 2010.

E. Carson and I. Daužickaitė. The stability of split-preconditioned FGMRES in four precisions. ETNA 2024.

E. Carson and I. Daužickaitė. Mixed precision sketching for least-squares problems and its application in GMRES-based iterative refinement. arXiv preprint arXiv:2410.06319, 2024.

### Goal

Still no comprehensive understanding of GMRES in mixed precision:

- How should we set  $u_g$ ,  $u_m$ , and  $u_a$ ?
- Which preconditioning style (left, right, flexible) should we use?
- In short,

# What are (*all*) the meaningful mixed precision GMRES strategies?

• *Meaningful* refers to a strategy for which lowering any of its precisions also lowers the bound on its attainable accuracy

- The numerical behavior of GMRES is characterized by two main properties: attainable accuracy and convergence rate
- Unfortunately, no strong theoretical results for the convergence rate (even in exact arithmetic)
  A. Greenbaum, V. Pták and Z. Strakoš. Any nonincreasing convergence curve is possible for GMRES. SIMAX 1996.
- $\Rightarrow$  We focus on the attainable accuracy

# Restarted GMRES as iterative refinement

### Restarted GMRES

1:  $x = x_0$ 2: r = b - Ax3: repeat 4:  $s_0 = M^{-1}r_0$ 5:  $\beta = ||s_0||, v_1 = s_0/\beta, k = 1$ 6: 7: 8: 9: for k = 1, ..., m do  $z_k = A v_k$  $w_k = M^{-1} z_k$ for i = 1, ..., k do 10:  $h_{i,k} = v_i^T w_k$ 11:  $w_k = w_k - h_{i,k}v_i$ 12: end for 13:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 14:  $V_k = [v_1, \ldots, v_k]$ 15:  $H_k = \{h_{i,i}\}_{1 \le i \le j+1; 1 \le j \le k}$ 16:  $y_k = \operatorname{argmin}_{Y} \|\beta e_1 - H_k y\|$ 17: end for 18:  $x = x + V_k y_k$ 19: r = b - Ax20: until  $||r|| \le \tau$ 

# Restarted GMRES as iterative refinement

### Restarted GMRES

1:  $x = x_0$ 2: r = b - Ax3: repeat 4:  $s_0 = M^{-1}r_0$ 5:  $\beta = \|s_0\|, v_1 = s_0/\beta, k = 1$ 6: 7: 8: for k = 1, ..., m do  $z_k = Av_k$  $w_k = M^{-1} z_k$ 9: for i = 1, ..., k do 10:  $h_{i,k} = v_i^T w_k$ 11:  $w_k = w_k - h_i k v_i$ 12: end for 13:  $h_{k+1,k} = ||w_k||, v_{k+1} = w_k/h_{k+1,k}$ 14:  $V_{k} = [v_1, \ldots, v_{k}]$ 15:  $H_k = \{h_{i,i}\}_{1 \le i \le j+1; 1 \le j \le k}$ 16:  $y_k = \operatorname{argmin}_v \|\beta e_1 - H_k y\|$ 17: end for 18:  $x = x + V_k y_k$ 19: r = b - Ax20: until  $||r|| \le \tau$ 

#### **Iterative refinement**

- 1:  $x = x_0$ 2: r = b - Ax3: repeat 4: Solve Ad = r5: x = x + d6: r = b - Ax7: until  $||r|| \le \tau$
- Restarted GMRES ⇔ iterative refinement with m iterations of GMRES as correction solver
- Attainable accuracy of IR:  $cond(A, x)u_r + u$
- Unlike GMRES, we also have a good understanding of IR's convergence rate:  $\frac{\|\hat{d}-d\|}{\|d\|}$
- ⇒ GMRES's attainable accuracy is IR's (and thus restarted GMRES's) convergence rate!

# Modular framework for the error analysis of GMRES

#### Error model and assumptions

We model GMRES as the following process:

We also make the following assumptions:

- $\label{eq:alpha} \begin{array}{l} \textbf{ 6} \\ \mbox{ All accuracy parameters are sufficiently less } \\ \mbox{ than 1: } \\ \mbox{ 0} \leq \varepsilon_c, \varepsilon_b, \varepsilon_{\rm ls}, \varepsilon_{\rm x} \ll 1. \end{array}$
- **6**  $Z_k$  is not numerically singular to accuracy  $\varepsilon_x$ :  $\kappa(Z_k)\varepsilon_x \ll 1$ .

#### Bound on the attainable accuracies

At the key iteration k, 
$$\hat{x}_k$$
 satisfies  $\frac{\|M_L^{-1}b - M_L^{-1}A\hat{x}_k\|}{\|M_L^{-1}A\|\|\hat{x}_k\| + \|M_L^{-1}b\|} \lesssim c(n,k)\xi$  and  $\frac{\|\hat{x}_k - x\|}{\|x\|} \lesssim c(n,k)\kappa(M_L^{-1}A)\xi$ ,  
where  $\xi = \alpha \varepsilon_c + \beta \varepsilon_b + \beta \varepsilon_{1s} + \lambda \varepsilon_x$ ,  $\alpha = \frac{\|M_L^{-1}AZ_k\|}{\|M_L^{-1}A\|\sigma_{\min}(Z_k)}$ ,  $\beta = \max(1, \alpha)$ ,  $\lambda = \kappa(Z_k)$ 

# Recovering known results with the framework

- Unpreconditioned Householder-GMRES is backward stable, i.e., ξ = O(u)
  I. Drkošová, A. Greenbaum, M. Rozložník, Z. Strakoš, Numerical stability of GMRES, BIT 1995.
- Unpreconditioned MGS-GMRES is backward stable, i.e., ξ = O(u)
  C. C. Paige, M. Rozložník, and Z. Strakoš. Modified Gram-Schmidt (MGS), least squares, and

backward stability of MGS-GMRES. SIMAX 2006.

- Flexible GMRES achieves  $\xi = O(u\kappa(Z_k))$ , hence:
  - FGMRES is not backward stable for general preconditioners
  - For a low precision LU preconditioner, under suitable assumptions, FGMRES is backward stable

B M. Arioli and I. S. Duff. Using FGMRES to obtain backward stability in mixed precision. ETNA 2008.

• Left and right preconditioned GMRES are also *not* backward stable in general:  $\xi = O(\frac{\|M^{-1}\|\|M\|}{\|M^{-1}A\|}u) \le O(\kappa(M)u)$ 

P. R. Amestoy, A. Buttari, N. J. Higham, J.-Y. L'Excellent, T. M., and B. Vieublé. Five-Precision GMRES-Based Iterative Refinement. SIMAX 2024.

### Deriving new results with the framework

- Unpreconditioned CGS2-GMRES is backward stable, i.e., ξ = O(u)
  A. Buttari, N. J. Higham, T. M., and B. Vieublé. A modular framework for the backward error analysis of GMRES. IMAJNA 2025.
- s-step (communication-avoiding) GMRES achieves ξ = O(uκ(Z<sub>k</sub>))
  E. Carson and Y. Ma. On the backward stability of s-step GMRES. arXiv preprint arXiv:2409.03079, 2024.
- Sketched GMRES also achieves ξ = O(uκ(Z<sub>k</sub>))
  L. Burke, E. Carson and Y. Ma. On the numerical stability of sketched GMRES. arXiv preprint arXiv:2503.19086, 2025.
- Mixed precision preconditioned GMRES

A. Buttari, X. Liu, T. M., and B. Vieublé. Mixed precision strategies for preconditioned GMRES: a comprehensive analysis. HAL EPrint hal-05071696, 2025.

# Applying the framework to mixed precision GMRES

We obtain the following bounds on the attainable forward error  $\frac{\|\hat{x}_k - x\|}{\|x\|}$ :

Left:

$$\begin{split} \kappa(\boldsymbol{M}^{-1}\boldsymbol{A})\boldsymbol{u_g} + \max(\rho,\kappa(\boldsymbol{M}^{-1}\boldsymbol{A}))\boldsymbol{u_m} + \kappa(\boldsymbol{A})\boldsymbol{u_a} \\ (\rho \leq \kappa(\boldsymbol{M}^{-1}\boldsymbol{A})\kappa(\boldsymbol{M}) \|\boldsymbol{A}\boldsymbol{v_j}\|/\|\boldsymbol{A}\|) \end{split}$$

• Right:

$$\kappa(AM^{-1})\kappa(M)u_{g} + \kappa(M)u_{m} + \kappa(A)u_{a}$$

• Flexible:

$$\kappa(AM^{-1})\kappa(M)u_{g} + \kappa(A)u_{a}$$

### Key observations

- For a given preconditioning style, the term in front of each precision is different
- For a given precision, the term in front of it depends on the preconditioning style

# Mixed precision strategies

	Left	Right	Flexible
$u_a = u_m \ll u_g$	exists	new	new
$u_a = u_g \ll u_m$	new	new	exists
$u_a \ll u_g = u_m$	new	new	new
$u_a \ll u_g \ll u_m$	new	new	new
$u_a \ll u_m \ll u_g$	new	new	new
$u_g \ll u_a = u_m$	new	new	new
$u_g \ll u_a \ll u_m$	new	new	new
$u_m \ll u_a = u_g$	new	new	new
$u_m \ll u_a \ll u_g$	new	new	new
$u_g = u_m \ll u_a$	new	new	new
$u_g \ll u_m \ll u_a$	new	new	new
$u_m \ll u_g \ll u_a$	new	new	new

# Mixed precision strategies

	Left	Right	Flexible
$u_a = u_m \ll u_g$	exists	new	new
$u_a = u_g \ll u_m$	new	new	exists
$u_a \ll u_g = u_m$	new	new	new
$u_a \ll u_g \ll u_m$	new	new	new
$u_a \ll u_m \ll u_g$	new	new	new
$u_g \ll u_a = u_m$	new	new	new
$u_g \ll u_a \ll u_m$	new	new	new
$u_m \ll u_a = u_g$	new	new	new
$u_m \ll u_a \ll u_g$	new	new	new
$u_g = u_m \ll u_a$	new	new	new
$u_g \ll u_m \ll u_a$	new	new	new
$u_m \ll u_g \ll u_a$	new	new	new

- Three floating-point arithmetics : fp64 (D), fp32 (S), and bfloat16 (H)
- Restarted GMRES with  $u_r = u = D$ , with reasonable tuning of the restart criterion. We report the total iteration count required to achieve

$$\frac{\|\widehat{x} - x\|_2}{\|x\|_2} \le 10^{-10}$$

Various matrices and preconditioners. A few representatives will be shown here; see paper for the full range of experiments:
 A. Buttari, X. Liu, T. M., and B. Vieublé. Mixed precision strategies for preconditioned GMRES: a comprehensive analysis. HAL EPrint hal-05071696, 2025.

1138_bus matrix, ILU preconditioner					
Iteration count					
Ua	Иg	u <sub>m</sub>	Left	Right	Flexible
D	D	D	17	18	18
S	S	S	34	37	38

1138_bus matrix, ILU preconditioner					
Iteration count					
u <sub>a</sub>	Иg	u <sub>m</sub>	Left	Right	Flexible
D	D	D	17	18	18
S	S	S	34	37	38
D	S	D	26	25	31

1138_bus matrix, ILU preconditioner						
		Iteration count				
Ua	Ug	u <sub>m</sub>	Left	Right	Flexible	
D	D	D	17	18	18	
S	S	S	34	37	38	
D	S	D	26	25	31	
D	D	S	25	28	18	

1138_bus matrix, ILU preconditioner						
	Iteration count					
u <sub>a</sub>	Иg	u <sub>m</sub>	Left	Right	Flexible	
D	D	D	17	18	18	
S	S	S	34	37	38	
D	S	D	<b>26</b>	25	31	
D	D	S	<b>25</b>	28	18	
D	S	S	27	31	31	

bc	bcsstk19 matrix, LU preconditioner				
Iteration count					
u <sub>a</sub>	ug	u <sub>m</sub>	Left	Right	Flexible
D	D	D	9	10	10
S	S	S	19	25	26
Н	Н	Н			
D	Н	Н	_		
S	Н	S	21		
D	Н	S	14	23	17
D	S	S	14	23	17
D	Н	D	13		

bwm200 matrix, LU preconditioner					
			Iteration count		
И <sub>а</sub>	u <sub>g</sub>	u <sub>m</sub>	Left	Right	Flexible
D	D	D	24	24	24
S	S	S	39	49	45
Н	Н	Н	274	729	—
S	S	Н	216	260	57
D	Н	Н	277	389	279
D	S	Н	216	250	44
D	D	Н	183	250	29
D	S	S	38	39	39

### Caveats

- The comparison between strategies is based on the *bounds* on their attainable accuracies, which are not necessarily sharp in practice
- The quantities appearing in these bounds are not always known or easy to compute
- The attainable accuracy of GMRES only becomes descriptive of the convergence rate of restarted GMRES for large restart sizes
  Y. Zhao, T. Fukaya, L. Zhang, and T. Iwashita. Numerical Investigation into the Mixed Precision GMRES(*m*) Method Using FP64 and FP32. Journal of Information Processing, 2022.
- Precisions are not allowed to vary across iterations

L. Giraud, S. Gratton, and J. Langou. Convergence in backward error of relaxed GMRES. SISC 2007.
 S. Gratton, E. Simon, D. Titley-Peloquin, P. L. Toint. A note on inexact inner products in GMRES. SIMAX 2022.

### Guidelines

 If your preconditioner is costly to apply: use u<sub>a</sub> = u<sub>g</sub> ≪ u<sub>m</sub> with flexible GMRES... ... or use a memory accessor with u<sub>a</sub> = u<sub>g</sub> = u<sub>m</sub>
 P. R. Amestoy, A. Jego, J.-Y. L'Excellent, T.M., and G. Pichon. BLAS-based Block Memory Accessor with Applications to Mixed Precision Sparse Direct Solvers. HAL EPrint hal-05019106, 2025.

- If the orthonormalization is expensive / if a large restart size is required: use  $u_a = u_m \ll u_g$  with left-preconditioned GMRES.
- If both of the above:

use  $u_a \ll u_m = u_g$  with left-preconditioned GMRES

## Conclusions

- Preconditioned GMRES is not backward stable (due to cancellation in the preconditioned matrix-vector product)
- Restarted GMRES is more accurate than GMRES (due to its iterative refinement properties)
- The precisions of the matrix-vector product, preconditioner application, and orthonormalization play very different roles
- Left, right, and flexible preconditioning are all very different, with no clear winner. Each tends to favor using low precision for different operations
- ⇒ Many (new) mixed precision strategies, with various levels of performance–accuracy tradeoff

# Thanks! Questions?