# Complexity and performance of the Block Low-Rank multifrontal factorization and its variants 

P. Amestoy*,1 A. Buttari*,2 J.-Y. L'Excellent ${ }^{\dagger, 3} \quad \underline{\text { T. Mary }}{ }^{*, 4}$

* Université de Toulouse ${ }^{\dagger}$ ENS Lyon
${ }^{1}$ INPT-IRIT ${ }^{2}$ CNRS-IRIT ${ }^{2}$ INRIA-LIP ${ }^{4}$ UPS-IRIT

SIAM PP'16, Paris Apr. 12-15

## Introduction

## Multifrontal (Duff '83) with Nested Dissection (George '73)



## Multifrontal (Duff '83) with Nested Dissection (George '73)



3D problem cost $\propto$
$\rightarrow$ Flops: $O\left(n^{2}\right)$, mem: $O\left(n^{4 / 3}\right)$


## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix

## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix

## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix

A block $B$ represents the interaction between two subdomains. If they have a small diameter and are far away their interaction is weak $\Rightarrow$ rank is low.

## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix
$A$ block $B$ represents the interaction between two subdomains. If they have a small diameter and are far away their interaction is weak $\Rightarrow$ rank is low.

$$
\tilde{B}=X Y^{\top} \text { such that } \operatorname{rank}(\tilde{B})=k_{\varepsilon} \text { and }\|B-\tilde{B}\| \leq \varepsilon
$$

If $k_{\varepsilon} \ll \operatorname{size}(B) \Rightarrow$ memory and flops can be reduced with a controlled loss of accuracy ( $\leq \varepsilon$ )

## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix

A block $B$ represents the interaction between two subdomains. If they have a small diameter and are far away their interaction is weak $\Rightarrow$ rank is low.

$$
\tilde{B}=X Y^{\top} \text { such that } \operatorname{rank}(\tilde{B})=k_{\varepsilon} \text { and }\|B-\tilde{B}\| \leq \varepsilon
$$

If $k_{\varepsilon} \ll \operatorname{size}(B) \Rightarrow$ memory and flops can be reduced with a controlled loss of accuracy ( $\leq \varepsilon$ )

## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix

- Leads to very low theoretical complexity
- Complex, hierarchical structure
- Simple structure
- Theoretical complexity?


## $\mathcal{H}$ and BLR matrices


$\mathcal{H}$-matrix


BLR matrix
$\Rightarrow$ Our hope is to find a good comprise between theoretical complexity and performance/usability

## Questions that will be answered in this talk

- Is the complexity of the BLR factorization asymptotically better than the full-rank one? (i.e., in $O\left(n^{\alpha}\right)$, with $\alpha<2$ and where $n$ is the number of unknowns)
- What are the different variants of the BLR factorization? Do they improve its complexity?
- How well does the complexity improvement translate into a performance gain?
- How parallel is the BLR factorization? What about its variants?

Variants of the BLR factorization

## Variants of the BLR LU factorization



- FSCU


## Variants of the BLR LU factorization



- FSCU (Factor,


## Variants of the BLR LU factorization



- FSCU (Factor, Solve,


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress,


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization


$\square$

- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization


$\square$
$\square$


- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization


$\square$
$\square$

- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve
- With LUA, no need to decompress accumulators


## Variants of the BLR LU factorization



- FSCU (Factor, Solve, Compress, Update)
- FSCU+LUA
- More natural in Left-looking
- Better granularity in update operations
- Potential recompression
- FCSU(+LUA)
- Restricted pivoting, e.g. to diagonal blocks
- Low-rank Solve
- Better ratio BLAS-3/BLAS-2 in Solve
- With LUA, no need to decompress accumulators

Complexity of the BLR factorization

- Extended theoretical work on $\mathcal{H}$-matrices by Hackbush and Bebendorf $(2003)$ and Bebendorf $(2005,2007)$ to the BLR case. Proof and computation of the theoretical complexity are available in On the Complexity of the Block Low-Rank Multifrontal Factorization, P. Amestoy, A. Buttari, J.-Y. L'Excellent and T. Mary (in preparation)
- Today, regarding the complexity, we focus on:
- Presenting some important properties of the BLR complexity
- Validating these properties experimentally


## Complexity of multifrontal BLR factorization

|  | operations (OPC) |  | factor size (NNZ) |  |
| :--- | :--- | :--- | :--- | :--- |
|  | $r=O(1)$ | $r=O\left(n^{\frac{1}{3}}\right)$ | $r=O(1)$ | $r=O\left(n^{\frac{1}{3}}\right)$ |
| FR | $O\left(n^{2}\right)$ | $O\left(n^{2}\right)$ | $O\left(n^{\frac{4}{3}}\right)$ | $O\left(n^{\frac{4}{3}}\right)$ |
| BLR FSCU | $O\left(n^{\frac{5}{3}}\right)$ | $O\left(n^{\frac{11}{6}}\right)$ | $O(n \log n)$ | $O\left(n^{\frac{4}{3}}\right)$ |
| BLR FSCU+LUA | $O\left(n^{\frac{14}{9}}\right)$ | $O\left(n^{\frac{16}{9}}\right)$ | $O(n \log n)$ | $O\left(n^{\frac{4}{3}}\right)$ |
| BLR FCSU+LUA | $O\left(n^{\frac{4}{3}}\right)$ | $O\left(n^{\frac{5}{3}} \log n\right)$ | $O(n \log n)$ | $O\left(n^{\frac{4}{3}}\right)$ |
| $\mathcal{H}$ | $O\left(n^{\frac{4}{3}}\right)$ | $O\left(n^{\frac{5}{3}}\right)$ | $O(n)$ | $O\left(n^{\frac{7}{6}}\right)$ |
| $\mathcal{H}$ (fully struct.) | $O(n)$ | $O\left(n^{\frac{4}{3}}\right)$ | $O(n)$ | $O\left(n^{\frac{7}{6}}\right)$ |

in the 3D case (similar analysis possible for 2D)
Important properties:

- The complexity of the standard BLR variant (FSCU) has a lower exponent than the full-rank one
- Each variant further improves the complexity, with the best one (FCSU+LUA) being not so far from the $\mathcal{H}$-case
- These properties hold for different rank bound assumptions, e.g. $r=O(1)$ or $r=O(N)=O\left(n^{\frac{1}{3}}\right)$

1. Poisson: $N^{3}$ grid with a 7 -point stencil with $u=1$ on the boundary $\partial \Omega$

$$
\Delta u=f
$$

2. Helmholtz: $N^{3}$ grid with a 27-point stencil, $\omega$ is the angular frequency, $v(x)$ is the seismic velocity field, and $u(x, \omega)$ is the time-harmonic wavefield solution to the forcing term $s(x, \omega)$.

$$
\left(-\Delta-\frac{\omega^{2}}{v(x)^{2}}\right) u(x, \omega)=s(x, \omega)
$$

$\omega$ is fixed and equal to 4 Hz .

## Experimental MF complexity: operations

OPC (Poisson, $\varepsilon=10^{-10}$ )


OPC (Helmholtz, $\varepsilon=10^{-5}$ )


- good agreement with theoretical complexity


## Experimental MF complexity: operations

OPC (Poisson, $\varepsilon=10^{-6}$ )


OPC (Helmholtz, $\varepsilon=10^{-4}$ )


- good agreement with theoretical complexity
- $\varepsilon$ only plays a role in the constant factor

Performance results

1. Distributed memory experiments are done on the eos supercomputer at the CALMIP center of Toulouse (grant 2014-P0989):

- Two Intel(r) 10-cores Ivy Bridge @ 2,8 GHz
- Peak per core is 22.4 GF/s
- 64 GB memory per node
- Infiniband FDR interconnect

2. Shared memory experiments are done on grunch at the LIP laboratory of Lyon:

- Two Intel(r) 14-cores Haswell @ 2,3 GHz
- Peak per core is $36.8 \mathrm{GF} / \mathrm{s}$
- Total memory is 768 GB

Poisson ( $\varepsilon=10^{-6}, N=192$ )


Helmholtz ( $\varepsilon=10^{-4}, N=192$ )


MPI+OpenMP parallelism (10 threads/MPI process, $1 \mathrm{MPI} /$ node)

- each time the number of processes doubles, speedup of $\sim 1.6$
- both FR and BLR scale reasonably well
- gain due to BLR remains constant


## Gains due to BLR (distributed, MPI+OpenMP)

Poisson $\left(\varepsilon=10^{-6}\right)$


Helmholtz $\left(\varepsilon=10^{-4}\right)$


- gains increase with problem size
- gain in flops does not fully translate into gain in time
- multithreaded efficiency lower in LR than in FR
- same remarks apply to Helmoltz, to a lesser extent


## Gains due to BLR (distributed, MPI+OpenMP)

Poisson $\left(\varepsilon=10^{-6}\right)$


Helmholtz $\left(\varepsilon=10^{-4}\right)$


- gains increase with problem size
- gain in flops does not fully translate into gain in time
- multithreaded efficiency lower in LR than in FR
- same remarks apply to Helmoltz, to a lesser extent
$\Rightarrow$ improve multithreading with variants

Focus on the Update step (which includes the Decompress)

|  |  | 1 thread |  | 28 threads |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | RL | LL | RL | LL |
| Poisson | FR | 62294 s | 65208 s | 3772 s | 4092 s |
| $(\mathrm{~N}=256)$ | BLR | 2516 s | 1544 s | 662 s | 183 s |
| Helmholtz | FR |  |  | 9862 s | 10234 s |
| $(\mathrm{~N}=256)$ | BLR |  |  | 1694 s | 1435 s |

## Right Looking Vs. Left-Looking (shared)

Focus on the Update step (which includes the Decompress)

|  |  | 1 thread |  | 28 threads |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | RL | LL | RL | LL |
| Poisson | FR | 62294 s | 65208 s | 3772 s | 4092 s |
| $(\mathrm{~N}=256)$ | BLR | 2516 s | 1544 s | 662 s | 183 s |
| Helmholtz | FR |  |  | 9862 s | 10234 s |
| $(\mathrm{~N}=256)$ | BLR |  |  | 1694 s | 1435 s |



## Right Looking Vs. Left-Looking (shared)

Focus on the Update step (which includes the Decompress)

|  |  | 1 thread |  | 28 threads |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | RL | LL | RL | LL |
| Poisson | FR | 62294 s | 65208 s | 3772 s | 4092 s |
| $(N=256)$ | BLR | 2516 s | 1544 s | 662 s | 183 s |
| Helmholtz | FR |  |  | 9862 s | 10234 s |
| $(N=256)$ | BLR |  |  | 1694 s | 1435 s |



## Right Looking Vs. Left-Looking (shared)

Focus on the Update step (which includes the Decompress)

|  |  | 1 thread |  | 28 threads |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | RL | LL | RL | LL |
| Poisson | FR | 62294 s | 65208s | 3772 s | 4092 s |
| $(N=256)$ | BLR | 2516 s | 1544 s | 662 s | 183 s |
| Helmoltz | FR |  |  | 9862 s | 10234 s |
| $(N=256)$ | BLR |  |  | 1694 s | 1435 s |



- in RL: FR (green) block is accessed many times; LR (blue) blocks are accessed once
- in LL: FR (green) block is accessed once; LR (blue) blocks are accessed many times
$\Rightarrow$ lower volume of memory transfers (more critical in multithreaded)


## Right Looking Vs. Left-Looking (shared)

Focus on the Update step (which includes the Decompress)

|  |  | 1 thread |  | 28 threads |  |
| :--- | :--- | ---: | ---: | ---: | ---: |
|  |  | RL | LL | RL | LL |
| Poisson | FR | 62294 s | 65208s | 3772 s | 4092 s |
| $(\mathrm{~N}=256)$ | BLR | 2516 s | 1544s | 662 s | 183 s |
| Helmoltz | FR |  |  | 9862 s | 10234 s |
| $(\mathrm{~N}=256)$ | BLR |  |  | 1694 s | 1435 s |



- in RL: FR (green) block is accessed many times; LR (blue) blocks are accessed once
- in LL: FR (green) block is accessed once; LR (blue) blocks are accessed many times
$\Rightarrow$ lower volume of memory transfers (more critical in multithreaded)
$\Rightarrow$ the Decompress part (135s) remains the bottleneck of the Update (183s)


## Performance of LUA (shared, 28 threads)

Double precision (d) performance benchmark of Decompress



|  | Poisson ( $N=256$ ) |  |  | Helmholtz ( $N=256$ ) |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | LL | LUA | LUA <br> +Rec.* | LL | LUA | LUA <br> +Rec.* |
| Flops in Update ( $\times 10^{13}$ ) | 1.0 | 1.0 | 0.58 | 43 | 43 | 30 |
| Avg. decompress size | 3.8 | 27.1 | 12.7 | 31.3 | 264.2 | 136.8 |
| Time in Update | 183s | 87s | 110s | 1435s | 1304s | 1295s |
| \% of peak reached | 5\% | 11\% | 5\% | 59\% | 65\% | 45\% |

* All metrics include the Recompression overhead


## Performance of LUA (shared, 28 threads)

Double precision (d) performance benchmark of Decompress


|  | Poisson ( $N=256$ ) |  |  | Helmholtz ( $N=256$ ) |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | LL | LUA | LUA <br> +Rec.* | LL | LUA | LUA <br> +Rec.* |
| Flops in Update ( $\times 10^{13}$ ) | 1.0 | 1.0 | 0.58 | 43 | 43 | 30 |
| Avg. decompress size | 3.8 | 27.1 | 12.7 | 31.3 | 264.2 | 136.8 |
| Time in Update | 183s | 87s | 110s | 1435s | 1304s | 1295s |
| \% of peak reached | 5\% | 11\% | 5\% | 59\% | 65\% | 45\% |

* All metrics include the Recompression overhead


## Performance of LUA (shared, 28 threads)

Double precision (d) performance benchmark of Decompress



|  | Poisson ( $N=256$ ) |  |  | Helmholtz ( $N=256$ ) |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | LL | LUA | LUA <br> +Rec.* | LL | LUA | LUA <br> +Rec.* |
| Flops in Update ( $\times 10^{13}$ ) | 1.0 | 1.0 | 0.58 | 43 | 43 | 30 |
| Avg. decompress size | 3.8 | 27.1 | 12.7 | 31.3 | 264.2 | 136.8 |
| Time in Update | 183s | 87s | 110s | 1435s | 1304s | 1295s |
| \% of peak reached | 5\% | 11\% | 5\% | 59\% | 65\% | 45\% |

* All metrics include the Recompression overhead


## Performance of BLR variants (shared, 28 threads)

Poisson ( $\varepsilon=10^{-6}, N=256$ )


Helmholtz ( $\varepsilon=10^{-3}, N=256$ )


- Non-computational time ( $\sim 300$ s) is not included $\Rightarrow$ addressed in MPI by tree parallelism and in OpenMP by W. Sid-Lakhdar's PhD thesis work (2014)

Poisson ( $\varepsilon=10^{-6}, N=256$ )


Helmholtz ( $\varepsilon=10^{-3}, N=256$ )


- Non-computational time ( $\sim 300$ s) is not included $\Rightarrow$ addressed in MPI by tree parallelism and in OpenMP by W. Sid-Lakhdar's PhD thesis work (2014)
- FCSU: Factor+Solve greatly reduced

Poisson ( $\varepsilon=10^{-6}, N=256$ )


Helmholtz ( $\varepsilon=10^{-3}, N=256$ )


- Non-computational time ( $\sim 300$ s) is not included $\Rightarrow$ addressed in MPI by tree parallelism and in OpenMP by W. Sid-Lakhdar's PhD thesis work (2014)
- FCSU: Factor+Solve greatly reduced
- LL: Update reduced thanks to lower volume of communications

Poisson ( $\varepsilon=10^{-6}, N=256$ )


Helmholtz ( $\varepsilon=10^{-3}, N=256$ )


- Non-computational time ( $\sim 300$ s) is not included $\Rightarrow$ addressed in MPI by tree parallelism and in OpenMP by W. Sid-Lakhdar's PhD thesis work (2014)
- FCSU: Factor+Solve greatly reduced
- LL: Update reduced thanks to lower volume of communications
- LUA: Update (Decompress) reduced thanks to better granularities

Poisson ( $\varepsilon=10^{-6}, N=256$ )


Helmholtz ( $\varepsilon=10^{-3}, N=256$ )


- Non-computational time ( $\sim 300$ s) is not included $\Rightarrow$ addressed in MPI by tree parallelism and in OpenMP by W. Sid-Lakhdar's PhD thesis work (2014)
- FCSU: Factor+Solve greatly reduced
- LL: Update reduced thanks to lower volume of communications
- LUA: Update (Decompress) reduced thanks to better granularities
- Recompression: potential flop reduction not translated into a time gain yet


## Conclusion and perspectives

## Complexity results

- Theoretical complexity of the BLR (multifrontal) factorization is asymptotically better than FR
- Studied BLR variants to further reduce complexity by achieving higher compression
- Numerical experiments show experimental complexity in agreement with theoretical one


## Performance results

- BLR variants possess better properties (efficiency, granularity, volume of communications, number of operations) $\Rightarrow$ leads to a considerable speedup w.r.t. standard BLR variant...
- ...which itself achieves up to 4.7 (Poisson) and 2.7 (Helmholtz) speedup w.r.t. FR


## Perspectives

- Implementation and performance analysis of the BLR variants in distributed memory (MPI+OpenMP parallelism)
- Efficient strategies to recompress accumulators (cf. J. Anton's talk)
- Pivoting strategies compatible with the BLR variants
- Influence of the BLR variants on the accuracy of the factorization


## Perspectives

- Implementation and performance analysis of the BLR variants in distributed memory (MPI+OpenMP parallelism)
- Efficient strategies to recompress accumulators (cf. J. Anton's talk)
- Pivoting strategies compatible with the BLR variants
- Influence of the BLR variants on the accuracy of the factorization


## Acknowledgements

- CALMIP, BULL and LIP for providing access to the machines
- SEISCOPE for providing the Helmholtz Generator
- LSTC members for scientific discussions


## Thanks! Questions?

Backup Slides

## Accumulator recompression



## Accumulator recompression



## Accumulator recompression



## Accumulator recompression



## Accumulator recompression



## Accumulator recompression

\section*{| $C$ |  |
| :---: | :---: |
|  | $C$ |
|  | $Q^{T}$ |}



- Weight recompression on $\left\{C_{i}\right\}_{i}$
$\Rightarrow$ With absolute threshold $\varepsilon_{\text {, each }} C_{i}$ can be compressed separately
- Redundancy recompression on $\left\{Q_{i}\right\}_{i}$
$\Rightarrow$ Bigger recompression overhead, when is it worth it?


## Experimental MF complexity: entries in factor

NNZ (Poisson)


NNZ (Helmholtz)


- good agreement with theoretical complexity
- $\varepsilon$ only plays a role in the constant factor

