# Can we avoid rounding-error estimation in HPC codes and still get trustworthy results?

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- Increasing power of current computers
  - → accelerators: GPUs, TPUs, FPGAs, etc.
- Enable to solve more complex problems
  - $\rightarrow~$  Quantum field theory, supernova simulation, etc.
- A high number of floating-point operations performed
  - $\rightarrow~$  Each of them can lead to a rounding error

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 $\Rightarrow$  Numerical validation is crucial

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  - execution time overhead
  - development cost induced by the application of numerical validation methods to HPC codes

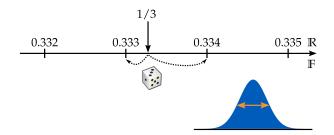
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Can we address this cost problem ...and still get trustworthy results?

Yes, when the input data is affected by rounding and/or measurement errors.

- Estimation of rounding errors: Discrete Stochastic Arithmetic (DSA) and the CADNA library
- error induced by perturbed data
- Our approach: combining DSA and standard floating-point arithmetic
- Numerical experiments
- Pros and cons of our approach

## Probabilistic approach for numerical validation



- operations are performed several times with random perturbations
  - → accuracy estimation
- analysis of the user code
  - → no specific numerical algorithms

Several tools:

CADNA [Chesneaux, 1990], MCAlib [Frechling et al., 2015], SAM [S. Graillat et al., 2011], VerifiCarlo [Denis et al., 2016], Verrou [Févotte et al., 2017]

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## Discrete Stochastic Arithmetic (DSA) [J. Vignes, 2004]

#### Principles

- each operation is executed 3 times with a random rounding mode:
   *R* → (*R*<sub>1</sub>, *R*<sub>2</sub>, *R*<sub>3</sub>) where each result *R<sub>i</sub>* is rounded up or down with the same probability
- the number of correct digits in the results is estimated using Student's test with the confidence level 95%
- operations are executed synchronously
  - ⇒ detection of numerical instabilities Ex: if (A>B) with A-B numerical noise
  - $\Rightarrow$  optimization of stopping criteria
    - Ex: stop when  $x_n x_{n+1}$  is numerical noise

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  - ⇒ optimization of stopping criteria Ex: stop when  $x_n - x_{n+1}$  is numerical noise

#### Implementations of DSA

- CADNA: for programs in double, single and/or half precision http://cadna.lip6.fr
- SAM: for arbitrary precision programs (based on MPFR) http://www-pequan.lip6.fr/~jezequel/SAM

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## The CADNA library http://cadna.lip6.fr



CADNA allows one to estimate rounding error propagation in any scientific program written in C, C++ or Fortran.

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CADNA provides new numerical types, the stochastic types, which consist of:

- 3 floating point variables
- an integer variable to store the accuracy.

All operators and mathematical functions are redefined for these types.

 $\Rightarrow$  CADNA requires only a few modifications in user programs.

Performance overhead:  $\times 4$  memory,  $\approx \times 10$  execution time

### An example without/with CADNA

```
Computation of P(x, y) = 9x^4 - y^4 + 2y^2 [S.M. Rump, 1983]
```

```
#include <stdio.h>
double rump(double x, double y) {
  return 9.0*x*x*x - v*v*v*v + 2.0*v*v:
}
int main(int argc, char **argv) {
  double x, y;
  x = 10864.0:
  v = 18817.0:
 printf("P1=%.14en", rump(x, y));
 x = 1.0/3.0:
 v = 2.0/3.0:
  printf("P2=%.14e\n". rump(x. v)):
 return 0:
}
```

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 v = 2.0/3.0:
  printf("P2=%.14e\n". rump(x. v)):
  return 0:
3
P1=2.00000000000000000e+00
```

```
P2=8.02469135802469e-01
```

```
#include <stdio.h>
```

```
double rump(double x, double y) {
  return 9.0*x*x*x-y*y*y*y+2.0*y*y;
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}
int main(int argc, char **argv) {
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    double x, y;
    x=10864.0; y=18817.0;
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 return 0;
}
```

```
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#include <cadna.h>
double_st rump(double_st x, double_st y) {
 return 9.0*x*x*x-y*y*y*y+2.0*y*y;
}
int main(int argc, char **argv) {
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```

only correct digits are displayed

Self-validation detection: ON Mathematical instabilities detection: ON Branching instabilities detection: ON Intrinsic instabilities detection: ON Cancellation instabilities detection: ON

P1= @.0 (no correct digits) P2= 0.802469135802469E+000

There are 2 numerical instabilities 2 LOSS(ES) OF ACCURACY DUE TO CANCELLATION(S)

A closer look at the floating-point values in P1 and P2:

P1=

-1.400000000000000e+01

-1.400000000000000e+01

2.000000000000000e+00

P2= 0.802469135802469e+00 0.802469135802469e+00

0.802469135802469e+00

#### Discrete Stochastic Arithmetic (DSA) and the CADNA library

#### Error induced by perturbed data

Our approach: combining DSA and standard floating-point arithmetic

4 Numerical experiments

5 Pros and cons of our approach

Let y = f(x) be an exact result and  $\hat{y} = \hat{f}(x)$  be the associated computed result.

- The forward error is the difference between y and  $\hat{y}$ .
- The backward analysis tries to seek for Δx s.t. ŷ = f(x + Δx).
   Δx is the backward error associated with ŷ.
   It measures the distance between the problem that is solved and the initial one.
- The condition number *C* of the problem is defined as:

$$C := \lim_{\varepsilon \to 0^+} \sup_{|\Delta x| \le \varepsilon} \left[ \frac{|f(x + \Delta x) - f(x)|}{|f(x)|} / \frac{|\Delta x|}{|x|} \right].$$

It measures the effect on the result of data perturbation.

## Error induced by perturbed data

The relative rounding error is denoted by u.

- *binary64* format (double precision):  $\mathbf{u} = 2^{-53}$
- *binary32* format (single precision):  $\mathbf{u} = 2^{-24}$ .

If the algorithm is backward-stable (*i.e.* the backward error is of the order of **u**)

 $|f(x) - \hat{f}(x)| / |f(x)| \lesssim C\mathbf{u}.$ 

If the input data are perturbed, *i.e.* the input data are not x but  $\hat{x} = x(1+\delta)$ , then one computes  $\hat{f}(\hat{x})$  with

$$|f(x) - \hat{f}(\hat{x})| / |f(x)| \lesssim C(\mathbf{u} + |\delta|).$$

If  $|\delta| \gg \mathbf{u}$ , the rounding error generated by  $\hat{f}$  is negligible w.r.t.  $C|\delta|$ .

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 $\Rightarrow$  Estimating this rounding error may be avoided.

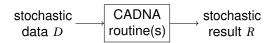
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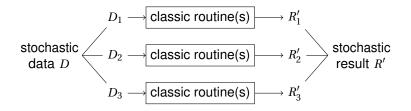
- 4 Numerical experiments
- 5 Pros and cons of our approach

- Computation routines are executed in a code that is controlled using DSA.
- Their input data are affected by errors (rounding errors and/or measurement errors).
- We compare 2 kinds of computation:
  - with a call to CADNA routines
  - with 3 calls to classic routines.



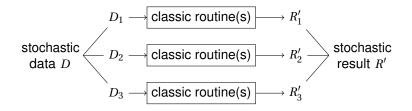
- D and R consist in stochastic arrays (each element is a triplet).
- Every arithmetic operation is performed 3 times with the random rounding mode.

## Our approach: computation with 3 calls to classic routines



- input data: 3 classic floating-point arrays  $D_1, D_2, D_3$  created from the triplets of D
- We get 3 classic floating-point arrays  $R'_1, R'_2, R'_3$ .
- A stochastic array *R*' created from *R*<sub>1</sub>', *R*<sub>2</sub>', *R*<sub>3</sub>' can be used in the next parts of the code.

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- $\Rightarrow$  we compare the number of correct digits (estimated by CADNA) in R and R'

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## Accuracy comparison

#### Data initialization

Each stochastic value is initialized as  $\alpha 10^e$ 

- $\alpha$ : random variable uniformly distributed in [-1,1]
- e: integer randomly generated in  $\{0, ..., E\}$  (DP: E = 20, SP: E = 3)
- $\Rightarrow$  generation of random data with different orders of magnitude.

#### Data perturbation

Each input value is perturbed with a **relative error**  $\delta$  using a CADNA function

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#### Accuracy analysis

For  $i = 1, ..., n^2$  (matrix mult.) or for i = 1, ..., n (matrix-vector mult.) we analyze:

- the accuracy  $C_{R^i}$  of the element  $R^i$  of R
- the accuracy  $C_{R'^i}$  of the element  $R'^i$  of R'

• 
$$\Delta^i = \left| C_{R^i} - C_{R'^i} \right|$$

## Accuracy comparison for matrix multiplication

Multiplication of square random matrices of size 500:

	accuracy		accuracy difference			
δ	of R		between R & R'			
	mean	min-max	mean	max		
double precision						
1.e-14	13.9	9-15	2.5e-02	2		
1.e-13	12.8	8-15	5.8e-03	1		
1.e-12	11.9	7-14	4.2e-04	1		
1.e-11	10.9	6-13	2.4e-05	1		
single precision						
1.e-6	5.6	1-7	2.3e-1	2		
1.e-5	4.8	0-7	1.9e-2	2		
1.e-4	3.7	0-6	2.8e-3	1		
1.e-3	2.8	0-5	2.8e-4	1		

- As the order of magnitude of  $\delta \nearrow$  the mean accuracy  $\searrow$  by 1 digit
- High perturbation in single precision  $\Rightarrow$  low accuracy on the results
- Low difference between the accuracy of R & R'

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## Accuracy comparison for matrix-vector multiplication

Multiplication of a square random matrix of size 1000 with a vector:

	accuracy		accuracy difference			
δ	of R		between R & R'			
	mean	min-max	mean	max		
double precision						
1.e-14	13.9	12-15	4.6e-02	1		
1.e-13	12.7	11-14	7.0e-03	1		
1.e-12	11.8	10-13	0	0		
1.e-11	10.9	9-12	0	0		
single precision						
1.e-6	5.5	3-7	3.2e-1	2		
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- As the order of magnitude of  $\delta \nearrow$  the mean accuracy  $\searrow$  by 1 digit
- High perturbation in single precision  $\Rightarrow$  low accuracy on the results
- The accuracy difference between *R* & *R'* remains low (in double precision, all the results have the same accuracy if  $\delta \ge 10^{-12}$ )

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Matrix and matrix-vector multiplication

We analyze the performance of various double precision codes.

### • "CADNA":

naive sequential multiplication with CADNA

"naive seq":

our approach using a sequential naive multiplication

• "naive OMP":

our approach using a naive parallel (OpenMP, 4 cores) multiplication

#### "MKL seq":

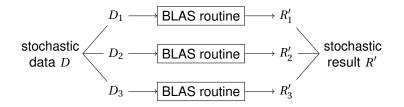
our approach using a sequential BLAS routine from the Intel MKL library

#### • "MKL OMP":

our approach using a parallel (OpenMP, 4 cores) MKL BLAS routine

Array copies except with CADNA

# Array copies in our experiments

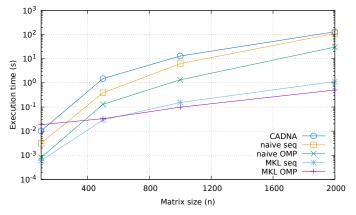


Conversions: array-of-structures ↔ structure-of-arrays

- before the BLAS routine: stochastic array → 3 classic arrays
- after the BLAS routine: 3 classic arrays  $\rightarrow$  stochastic array
- Worst case (maximum array copy cost in total execution time)
   BLAS routines continuously used
   ⇒ array copies only before and after them
- Both computation and array copies parallelized in the OpenMP codes

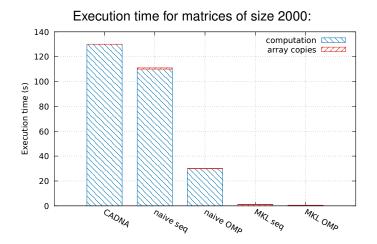
# Performance for matrix multiplication

Execution time including matrix multiplications and array copies:



- Despite memory copies, the codes using 3 classic matrix multiplications perform better than the CADNA routine.
- For matrices of size 2000, the MKL OpenMP implementation outperforms the CADNA routine by a factor 294.

## Performance for matrix multiplication



Most of the execution time is spent in matrix multiplication.

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# Performance for matrix multiplication

CADNA vs our approach with MKL OMP

#### Core i7-8650U (1.9 GHz, 4 cores), n=2000:

	CADNA	Proposed w/	Speedup
		MKL OMP	
Comp	130	0.393	331x
Сору	_	0.0495	—
Total	130	0.4425	294x

Dual-socket Xeon Gold 6126 (2.6 GHz, 12 cores×2), n=5000:

	CADNA	Proposed w/	Speedup
		MKL OMP	
Comp	2520	0.563	4476x
Сору	_	0.0889	—
Total	2520	0.652	3865x

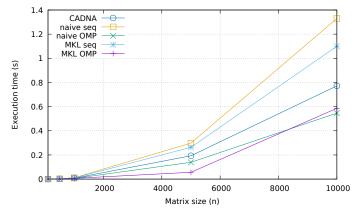
On large scale:

- the performance gain increases
- the array copy cost becomes visible

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## Performance for matrix-vector multiplication

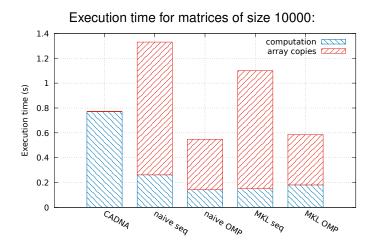
Execution time including matrix-vector multiplications and array copies:



- The CADNA routine performs better than the other sequential codes.
- From a certain matrix size, the OpenMP codes that use classic floating-point arithmetic perform better than the CADNA code.

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## Performance for matrix-vector multiplication



 In the sequential codes that use classic floating-point arithmetic the main part of the execution time is spent in array copies.

### Discrete Stochastic Arithmetic (DSA) and the CADNA library

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#### Pros

#### • performance gain:

- DSA operations are avoided
- use of vendor optimized libraries
- applicability:
  - no code translation to a CADNA version

### Cons

#### we lose CADNA features:

- instability detection
- accuracy improvement

# Instability detection

Without CADNA:

- numerical instabilities are not detected <sup>(C)</sup>
- ullet results with no correct digits appear as numerical noise igodot

#### Example: matrix multiplication with catastrophic cancellations

### Input data: square matrices A & B of size 10 in double precision

- 1st line of A: [1,...,1,-1,...,-1] (1st half: 1, 2nd half: -1)
- each element of B set to 1
- A and B pertubed with a relative error  $\delta = 10^{-12}$

### Results: C = A \* B with CADNA, C' = A \* B without CADNA

• 1st line of C and C': @.0 (numerical noise, triplet with no common digits)

With CADNA:

• 10 catastrophic cancellations are detected.

# Accuracy improvement with CADNA

#### Example: Gauss algorithm with pivoting

### Input data:

We solve in single precision the system Ax = b with

$$A = \begin{pmatrix} 21 & 130 & 0 & 2.1 \\ 13 & 80 & 4.74 & 10^8 & 752 \\ 0 & -0.4 & 3.9816 & 10^8 & 4.2 \\ 0 & 0 & 1.7 & 9 & 10^{-9} \end{pmatrix} \quad b = \begin{pmatrix} 153.1 \\ 849.74 \\ 7.7816 \\ 2.6 & 10^{-8} \end{pmatrix}$$

A and b perturbed with a relative error  $\delta = 10^{-6}$ 

### Results: x with CADNA, x' without CADNA

$$x = \begin{pmatrix} 0.100 \pm +001 \\ 0.999 \pm +000 \\ 0.999999 \pm -008 \\ 0.999999 \pm +000 \end{pmatrix} \quad x' = \begin{pmatrix} @.0 \\ @.0 \\ @.0 \\ 0.999999 \pm +000 \end{pmatrix} \quad x_{exact} = \begin{pmatrix} 1 \\ 1 \\ 10^{-8} \\ 1 \end{pmatrix}$$

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### Example: Gauss algorithm with pivoting

### Results: x with CADNA, x' without CADNA

$$x = \begin{pmatrix} 0.100 \text{E} + 001 \\ 0.999 \text{E} + 000 \\ 0.999999 \text{E} - 008 \\ 0.9999999 \text{E} + 000 \end{pmatrix} \quad x' = \begin{pmatrix} @.0 \\ @.0 \\ @.0 \\ 0.999999 \text{E} + 000 \end{pmatrix} \quad x_{exact} = \begin{pmatrix} 1 \\ 1 \\ 10^{-8} \\ 1 \end{pmatrix}$$

## Test for pivoting: if $(|A_{i,j}| > p_{max})$ ...

With CADNA a non-significant element is not chosen as a pivot.

### Instabilities detected by CADNA:

There are 3 numerical instabilities

- 1 UNSTABLE BRANCHING(S)
- 1 UNSTABLE INTRINSIC FUNCTION(S)
- 1 LOSS(ES) OF ACCURACY DUE TO CANCELLATION(S)

- In a code controlled using CADNA, if computation-intensive routines are run with perturbed data,
  - classic BLAS routines can be executed 3 times instead of the CADNA routines with almost no accuracy difference on the results
  - the performance gain can be high with BLAS routines from an optimized library
  - but we lose the instability detection.
- The same conclusions would be valid with an HPC code using MPI.
   In the same conditions (computation-intensive routines & perturbed data) CADNA-MPI routine ⇒ optimized floating-point MPI routines.
- Application of our approach to real-life examples with realistic data sets.

Thanks for your attention!