

# Computation of pseudozero abscissa

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SYNACS 2004

Timișoara, Romania, September 26 - 30, 2004



DALI

*Digital Architectures et Logiciels Informatiques*



# Motivations

Polynomial coefficients are often approximate values

Three well known sources of approximation are considered in scientific computation :

- (1) errors due to discretization and truncation,
- (2) errors due to roundoff, and
- (3) errors due to uncertainty in the data.

⇒ Use tools designed for such approximate polynomials in control theory

# Outline of the talk

## 1 — Pseudozero set

- Definition
- Computation

## 2 — Applications of pseudozeros in control theory

- Robust stability of polynomials
- Pseudozero abscissa of polynomials

# Pseudozeros : definition, computation

## Pseudozero set : definition

Let  $p$  be a given polynomial of  $\mathbf{C}_n[z]$

### Perturbation :

Neighborhood of polynomial  $p$

$$N_\varepsilon(p) = \{\hat{p} \in \mathbf{C}_n[z] : \|p - \hat{p}\| \leq \varepsilon\}.$$

### Definition of the $\varepsilon$ -pseudozero set :

$$Z_\varepsilon(p) = \{z \in \mathbb{C} : \hat{p}(z) = 0 \text{ for } \hat{p} \in N_\varepsilon(p)\}.$$

$\|\cdot\|$  a norm on the vector of the coefficients of  $p$

Pseudozero set : the set of the zeros of polynomials “near  $p$ ”.

# Pseudozeros are easily computable

## Theorem [Stetter] :

The  $\varepsilon$ -pseudozeros set satisfies

$$Z_\varepsilon(p) = \left\{ z \in \mathbb{C} : |g(z)| := \frac{|p(z)|}{\|\underline{z}\|_*} \leq \varepsilon \right\},$$

where  $\underline{z} = (1, z, \dots, z^n)$  and  $\|\cdot\|_*$  is the dual norm of  $\|\cdot\|$ ,

$$\|y\|_* = \sup_{x \neq 0} \frac{|y^* x|}{\|x\|}$$

# Pseudzero set : algorithm of computation

1. We mesh a square containing all the roots of  $p$  (MATLAB command : `meshgrid`).
2. We compute  $g(z) := \frac{|p(z)|}{\|z\|_*}$  for all the nodes  $z$  of the grid.
3. We plot the contour level  $|g(z)| = \varepsilon$  (MATLAB command : `contour`).

## Initialization :

- Find a square containing all the roots of  $p$  and all the pseudozeros.
- Find a grid step that separates all the roots.

## A famous example

Pseudozero set of the *Wilkinson* polynomial

$$\begin{aligned}W_{20} &= (z - 1)(z - 2) \cdots (z - 20), \\ &= z^{20} - 210z^{19} + \cdots + 20!.\end{aligned}$$

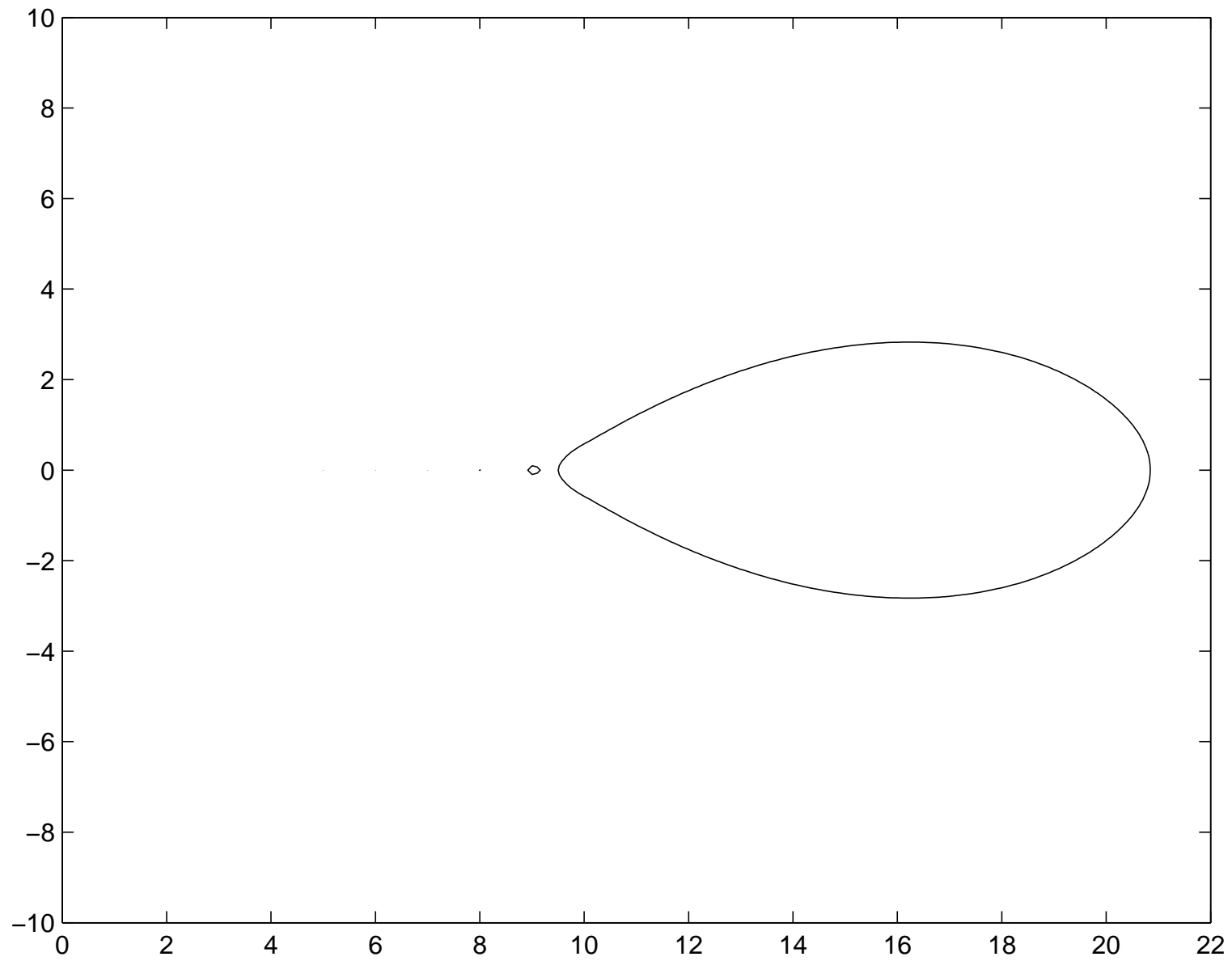
We only perturb the coefficient of  $z^{19}$  with  $\varepsilon = 2^{-23}$ .

One uses the weighted-norm  $\|\cdot\|_\infty$  :

$$\|p\|_\infty = \max_i \frac{|p_i|}{m_i} \text{ with } m_i \text{ non negative}$$

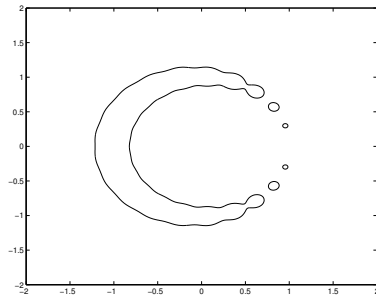
with  $m_{19} = 1$ ,  $m_i = 0$  otherwise and the convention  $m/0 = \infty$  if  $m > 0$  and  $0/0 = 0$ .



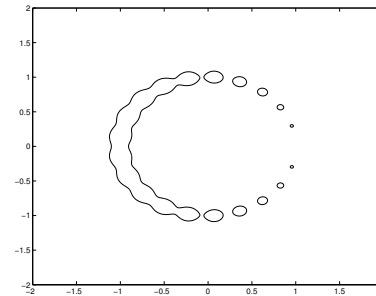


# Evolution of $\varepsilon$ -pseudozero w.r.t $\varepsilon$

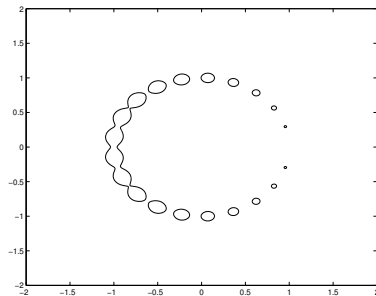
Pseudozero set of the polynomial  $p(z) = 1 + z + \dots + z^{20}$  for different values of  $\varepsilon$  (for the 2-norm).



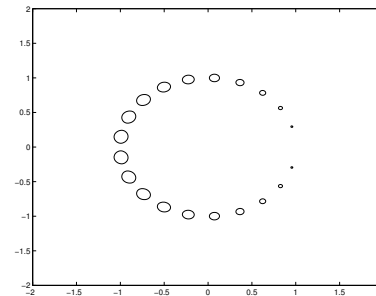
(a)  $\varepsilon = 10^{-1}$



(b)  $\varepsilon = 10^{-1.2}$



(c)  $\varepsilon = 10^{-1.3}$



(d)  $\varepsilon = 10^{-1.4}$

# Pseudozeros : brief survey of existing references

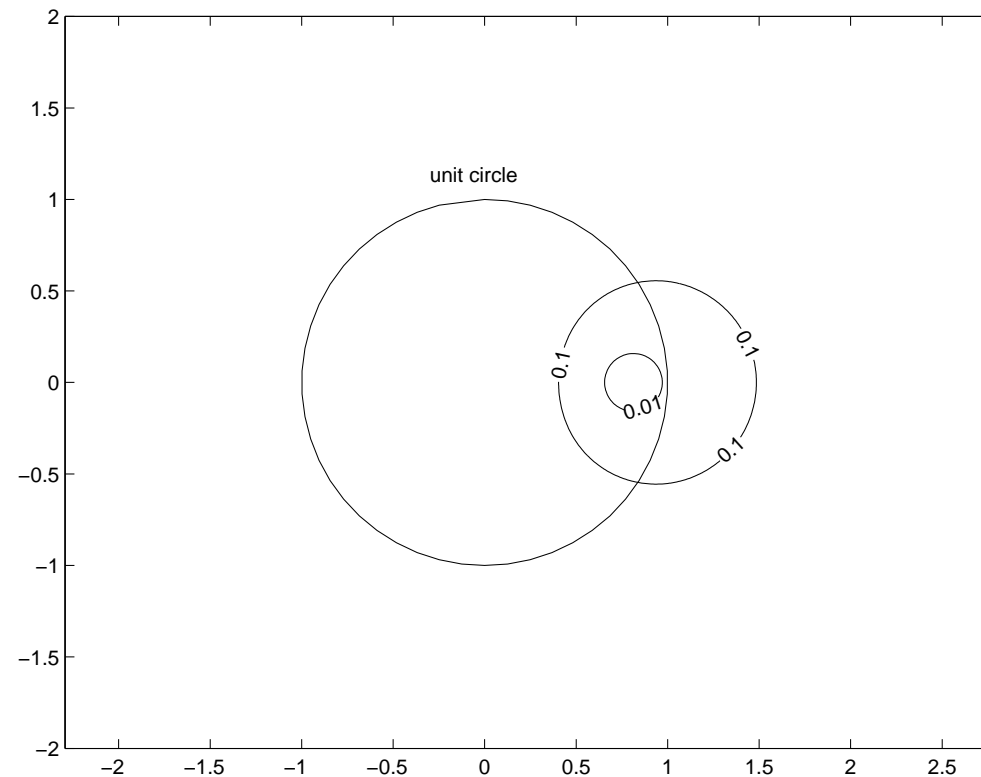
- ▶ Mosier (1986) : Definition and study for the  $\infty$ -norm.
- ▶ Hinrichsen and Kelb (1993) : Spectral value sets.
- ▶ Trefethen and Toh (1994) : Study for the 2-norm.  
pseudozeros  $\approx$  pseudospectra of the companion matrix.
- ▶ Zhang (2001) : Study the influence of the basis for the 2-norm (condition number of the evaluation).
- ▶ Stetter (2004) : *Numerical Polynomial Algebra* (SIAM). Generalization of the previous works.

# Robust stability and Pseudozero abscissa

# Schur robust stability in control theory

Schur stability : |roots of  $p$ | < 1.

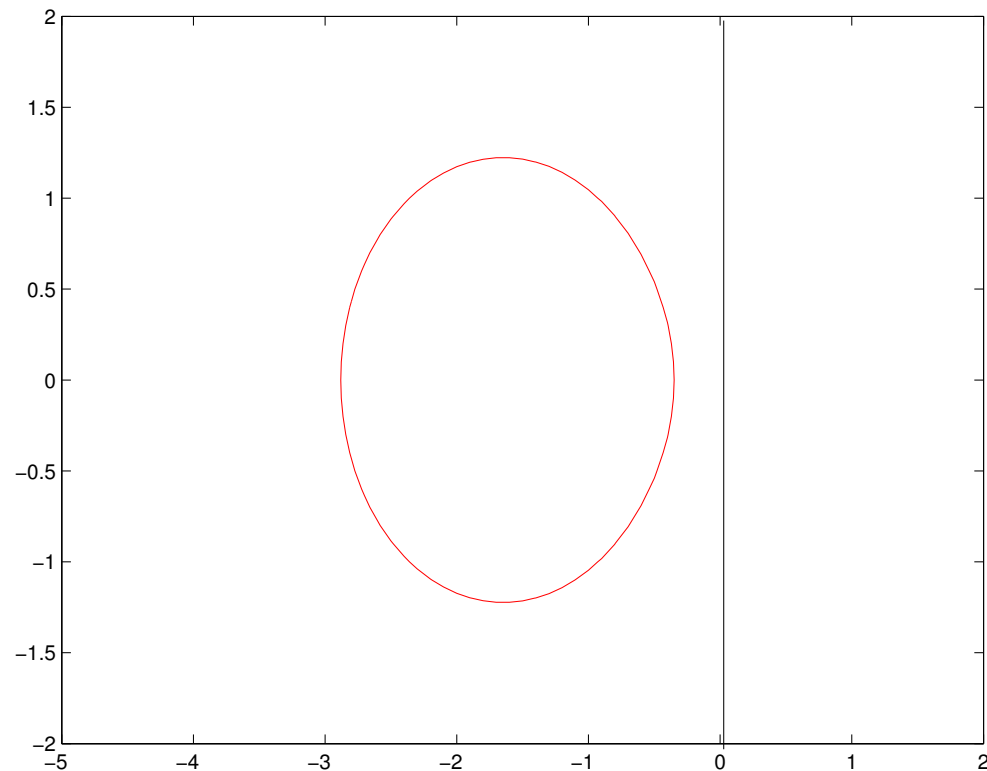
$\varepsilon$ -pseudozero set of  $p(z) = (z - 0.8)^2$  for  $\varepsilon = 0.1$  and  $\varepsilon = 0.01$ .



# Hurwitz robust stability in control theory

Hurwitz stability : Real part of roots of  $p < 0$ .

$\varepsilon$ -pseudozero set of  $p(z) = (z + 1)^2$  for  $\varepsilon = 0.4$ .



# Computation of pseudozero abscissa

$\mathcal{P}_n$  : polynomials of  $\mathbf{C}[X]$  of degree at most  $n$

$\mathcal{M}_n$  : monic polynomials of  $\mathcal{P}_n$  of degree  $n$

$\|\cdot\|$  : the 2-norm of the coefficients of a polynomial

**Definition.** A polynomial is **stable** if all its roots have negative real part and **unstable** otherwise (Hurwitz stability).

The function *abscissa*  $a : \mathcal{P} \rightarrow \mathbf{R}$  is defined by

$$a(p) = \max\{\operatorname{Re}(z) : p(z) = 0\}.$$

A polynomial  $p$  is stable  $\iff a(p) < 0$

# Motivation

In control theory, transfer functions are often written as  $H(p) = \frac{N(p)}{D(p)}$  where  $N$  and  $D$  are polynomials.

The system is stable if  $D$  is a stable polynomial .

Question : if  $D$  is stable, is it still stable when perturbed

(we assume that  $D$  is monic)



# Pseudzero abscissa mapping

$\varepsilon$ -pseudzero abscissa mapping  $a_\varepsilon : \mathcal{P}_n \rightarrow \mathbf{R}$  :

$$a_\varepsilon(p) = \max\{\operatorname{Re}(z) : z \in Z_\varepsilon(p)\}.$$

**Statement of the problem :**

Given a polynomial  $p \in \mathcal{M}_n$ , let us compute  $a_\varepsilon(p)$ .

A polynomial  $p$  is  $\varepsilon$ -robustly stable  $\iff a_\varepsilon(p) < 0$

# Our solutions

## Tools

- an explicit formula that defines the **pseudozeros**
- the **continuous dependency** of the roots w.r.t the polynomial **coefficients**
- **Sturm sequences** to count the real roots
- *criss-cross* algorithm

## The results : 3 algorithms

- a plotting algorithm
- a bisection algorithm
- a criss-cross algorithm

# Pseudozero set for monic polynomials

**Perturbation** : Neighborhood of polynomial  $p$

$$N_\varepsilon(p) = \{\hat{p} \in \mathcal{M}_n : \|p - \hat{p}\| \leq \varepsilon\}.$$

**Definition of the  $\varepsilon$ -pseudozero set :**

$$Z_\varepsilon(p) = \{z \in \mathbb{C} : \hat{p}(z) = 0 \text{ for } \hat{p} \in N_\varepsilon(p)\}.$$

$\|\cdot\|$  is the 2-norm on the vector of the coefficients of  $p$

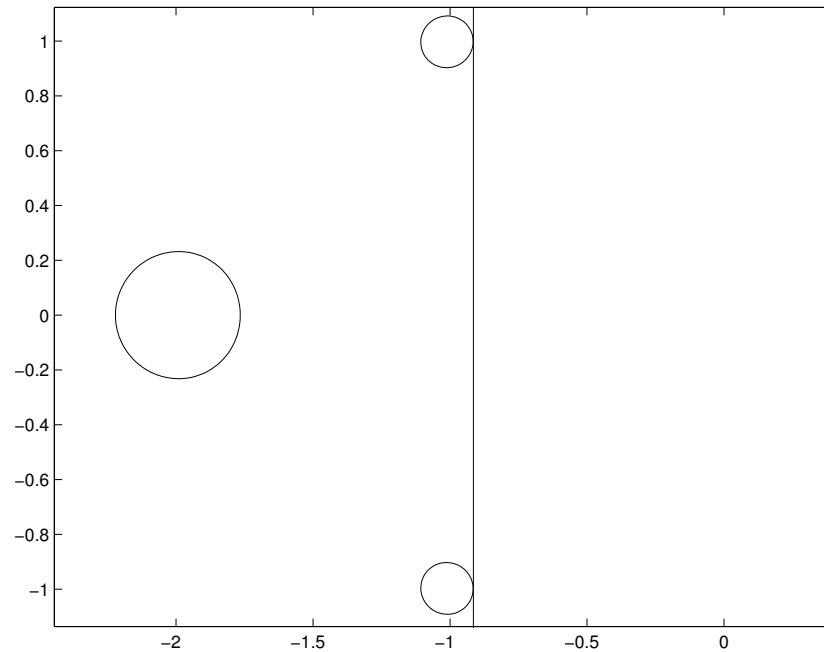
The  $\varepsilon$ -pseudozeros set satisfies

$$Z_\varepsilon(p) = \left\{ z \in \mathbb{C} : |g(z)| := \frac{|p(z)|}{\|\underline{z}\|} \leq \varepsilon \right\},$$

where  $\underline{z} = (1, z, \dots, z^{n-1})$

## A plotting algorithm

- Draw the  $\varepsilon$ -pseudozero set
- Draw the vertine line that intersects the right-most point within the  $\varepsilon$ -pseudozero set



$\varepsilon$ -pseudozero set of  $p(z) = z^3 + 4z^2 + 6z + 4$  for  $\varepsilon = 0.1$

$$a_\varepsilon(p) \approx -0.9$$

## Another characterization of $Z_\varepsilon(p)$

Let us denote  $h_{p,\varepsilon} : \mathbf{R}^2 \rightarrow \mathbf{R}$ , the function

$$h_{p,\varepsilon}(x, y) = |p(x + iy)|^2 - \varepsilon^2 \sum_{j=0}^{n-1} (x^2 + y^2)^j.$$

Then one has

$$Z_\varepsilon(p) = \{(x, y) \in \mathbf{R}^2 : h_{p,\varepsilon}(x, y) \leq 0\}$$

$\implies h_{p,\varepsilon}(\cdot, y)$  et  $h_{p,\varepsilon}(x, \cdot)$  are polynomials of degree  $2n$ .

**Theorem.** For any real  $x \geq a(p)$ , the equation  $h_{p,\varepsilon}(x, y) = 0$  has a real solution  $y$  if and only if  $x \leq a_\varepsilon(p)$ .

# A bisection algorithm

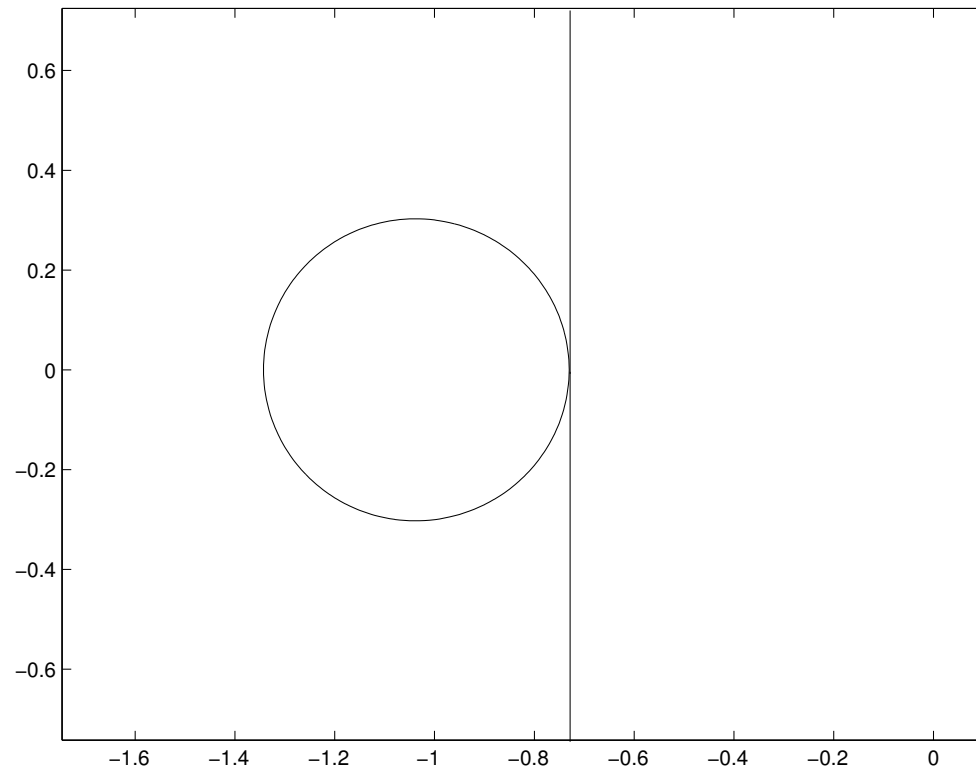
**Require :** a stable polynomial  $p$ , the parameter  $\varepsilon$  and a tolerance  $\tau$

**Ensure :** a number  $\alpha$  such that  $|\alpha - a_\varepsilon(p)| \leq \tau$

- 1:  $\gamma := a(p), \quad \delta := 1 + \|p\| + \varepsilon$
- 2: **while**  $|\gamma - \delta| > \tau$  **do**
- 3:      $x := \frac{\gamma + \delta}{2}$
- 4:     **if** the equation  $h_{p,\varepsilon}(x, y) = 0$  has a real solution **then**
- 5:          $\delta := x$
- 6:     **else**
- 7:          $\gamma := x$
- 8:     **end if**
- 9: **end while**
- 10: **return**  $\alpha = \frac{\gamma + \delta}{2}$

## Numerical simulation

For  $p(z) = z^5 + 5^4 + 10z^3 + 10z^2 + 5z + 1$ , the algorithm gives  $a_\varepsilon(p) \approx -0.719669$  for  $\varepsilon = 0.001$  and  $\tau = 0.00001$



$\varepsilon$ -pseudozero set of  $p(z) = z^5 + 5^4 + 10z^3 + 10z^2 + 5z + 1$  for  $\varepsilon = 0.001$

# A criss-cross algorithm

**Require** : a polynomial  $p$ , the parameter  $\varepsilon$

- 1: **Initialize** :  $x^1 = a(p)$  and  $r = 1$
- 2: **Vertical search** : find open intervals  $I_1^r, \dots, I_{l_r}^r$  where  $h(x^r, y) < 0$  for  $y \in \cup_{k=1}^{l_r} I_k^r$
- 3: **Horizontal search** : for each  $I_k^r$ , define  $\omega_k^r = \text{midpoint}(I_k^r)$  and find the largest real zeros  $x_k^r$  of the function  $h(\cdot, \omega_k^r)$  for  $k = 1 : l_r$
- 4: **Define**  $x^{r+1} = \max\{x_k^r, k = 1, \dots, l_r\}$ , increment  $r$  by one and return to Step 2.



# Conclusion and future work

## Conclusion :

Pseudozero set provides

- a better understanding of the effect of **coefficient perturbations**
- some applications for **robust stability**

## Future work :

- an analysis of the **convergence** of the **criss-cross algorithm** (we hope a quadratic convergence)
- an **implementation** of the criss-cross algorithm
- a generalization to **pseudospectra of polynomial matrices**